

# Modular Heap Shape Analysis for Java Programs

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**AProVE: termination and complexity analysis tool for Java**

## **AProVE: termination and complexity analysis tool for Java**

```
public void add(Object x) {  
    List l = this;  
    while (l.n != null) {  
        l = l.n;  
    }  
    List ll = new List();  
    l.n = ll;  
    ll.v = x;  
}
```

## AProVE: termination and complexity analysis tool for Java

```
public void add(Object x) {  
    List l = this;  
    while (l.n != null) {  
        l = l.n;  
    }  
    List ll = new List();  
    l.n = ll;  
    ll.v = x;  
}
```

- sophisticated heap shape analysis

## AProVE: termination and complexity analysis tool for Java

```
public void add(Object x) {  
    List l = this;  
    while (l.n != null) {  
        l = l.n;  
    }  
    List ll = new List();  
    l.n = ll;  
    ll.v = x;  
}
```

- sophisticated heap shape analysis
- lacks **modularity**

```
public void add(Object x) {  
    List l = this;  
  
    while (l.n != null) {  
        l = l.n;  
  
    }  
  
    List ll = new List();  
  
    l.n = ll;  
  
    ll.v = x;  
  
}
```

```
public void add(Object x) {  
   $\langle \text{this} = o_1, x = o_2 \mid \varepsilon \rangle$   
  List l = this;  
  
  while (l.n != null) {  
    l = l.n;  
  
  }  
  
  List ll = new List();  
  
  l.n = ll;  
  
  ll.v = x;  
  
}
```

```
public void add(Object x) {  
   $\langle \text{this} = o_1, x = o_2 \mid \varepsilon \rangle$   
  List l = this;  
   $\langle \text{this} = o_1, x = o_2, l = o_1 \mid \varepsilon \rangle$   
  
  while (l.n != null) {  
  
    l = l.n;  
  
  }  
  
  List ll = new List();  
  
  l.n = ll;  
  
  ll.v = x;  
  
}
```



```
public void add(Object x) {  
  <this = o1, x = o2 | ε>  
  List l = this;  
  <this = o1, x = o2, l = o1 | ε>  
  
  while (l.n != null) {  
    <this = o1, x = o2, l = o1 | ε>  
    l = l.n;  
  
  }  
  
  List ll = new List();  
  
  l.n = ll;  
  
  ll.v = x;  
  
}
```

```
public void add(Object x) {  
  ⟨this = o1, x = o2 | ε⟩  
  List l = this;  
  ⟨this = o1, x = o2, l = o1 | ε⟩  
  
  while (l.n != null) {  
    ⟨this = o1, x = o2, l = o1 | ε⟩  
    l = l.n;  
    ⟨this = o1, x = o2, l = o3 | o1  $\xrightarrow{n}$  o3⟩  
  }  
  
  List ll = new List();  
  
  l.n = ll;  
  
  ll.v = x;  
  
}
```

```
public void add(Object x) {  
   $\langle \text{this} = o_1, x = o_2 \mid \varepsilon \rangle$   
  List l = this;  
   $\langle \text{this} = o_1, x = o_2, l = o_1 \mid \varepsilon \rangle$   
   $\langle \text{this} = o_1, x = o_2, l = o_3 \mid o_1 = ? o_3, o_1 \not\sim o_3 \rangle$   
  while (l.n != null) {  
     $\langle \text{this} = o_1, x = o_2, l = o_1 \mid \varepsilon \rangle$   
    l = l.n;  
     $\langle \text{this} = o_1, x = o_2, l = o_3 \mid o_1 \xrightarrow{n} o_3 \rangle$   
  
  }  
  
  List ll = new List();  
  
  l.n = ll;  
  
  ll.v = x;  
  
}
```

```
public void add(Object x) {  
  ⟨this = o1, x = o2 | ε⟩  
  List l = this;  
  ⟨this = o1, x = o2, l = o1 | ε⟩  
  ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 ↘ o3⟩  
  while (l.n != null) {  
    ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 ↘ o3⟩  
    l = l.n;  
    ⟨this = o1, x = o2, l = o3 | o1  $\xrightarrow{n}$  o3⟩  
  }  
  
  List ll = new List();  
  
  l.n = ll;  
  
  ll.v = x;  
  
}
```

```

public void add(Object x) {
  ⟨this = o1, x = o2 | ε⟩
  List l = this;
  ⟨this = o1, x = o2, l = o1 | ε⟩
  ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 ↘ o3⟩
  while (l.n != null) {
    ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 ↘ o3⟩
    l = l.n;
    ⟨this = o1, x = o2, l = o4 | o1 =? o3, o1 ↘ o3, o3  $\xrightarrow{n}$  o4, o1 =? o4, o1 ↘ o4⟩
  }

  List ll = new List();

  ll.n = ll;

  ll.v = x;

}

```

```

public void add(Object x) {
  ⟨this = o1, x = o2 | ε⟩
  List l = this;
  ⟨this = o1, x = o2, l = o1 | ε⟩
  ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 ↘ o3⟩
  while (l.n != null) {
    ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 ↘ o3⟩
    l = l.n;
    ⟨this = o1, x = o2, l = o4 | o1 =? o3, o1 ↘ o3, o3 n→ o4, o1 =? o4, o1 ↘ o4⟩
    ⟨this = o1, x = o2, l = o4 | o1 =? o4, o1 ↘ o4⟩
  }

  List ll = new List();

  l.n = ll;

  ll.v = x;

}

```

```

public void add(Object x) {
  ⟨this = o1, x = o2 | ε⟩
  List l = this;
  ⟨this = o1, x = o2, l = o1 | ε⟩
  ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 ↘ o3⟩
  while (l.n != null) {
    ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 ↘ o3⟩
    l = l.n;
    ⟨this = o1, x = o2, l = o4 | o1 =? o3, o1 ↘ o3, o3 n→ o4, o1 =? o4, o1 ↘ o4⟩
    ⟨this = o1, x = o2, l = o4 | o1 =? o4, o1 ↘ o4⟩
  }
  ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 ↘ o3⟩
  List ll = new List();

  l.n = ll;

  ll.v = x;

}

```

```

public void add(Object x) {
  ⟨this = o1, x = o2 | ε⟩
  List l = this;
  ⟨this = o1, x = o2, l = o1 | ε⟩
  ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 ↘ o3⟩
  while (l.n != null) {
    ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 ↘ o3⟩
    l = l.n;
    ⟨this = o1, x = o2, l = o4 | o1 =? o3, o1 ↘ o3, o3 n→ o4, o1 =? o4, o1 ↘ o4⟩
    ⟨this = o1, x = o2, l = o4 | o1 =? o4, o1 ↘ o4⟩
  }
  ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 ↘ o3⟩
  List ll = new List();
  ⟨this = o1, x = o2, l = o3, ll = o5 | o1 =? o3, o1 ↘ o3⟩
  ll.n = ll;

  ll.v = x;
}

```



```

public void add(Object x) {
  ⟨this = o1, x = o2 | ε⟩
  List l = this;
  ⟨this = o1, x = o2, l = o1 | ε⟩
  ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 ↘ o3⟩
  while (l.n != null) {
    ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 ↘ o3⟩
    l = l.n;
    ⟨this = o1, x = o2, l = o4 | o1 =? o3, o1 ↘ o3, o3 n→ o4, o1 =? o4, o1 ↘ o4⟩
    ⟨this = o1, x = o2, l = o4 | o1 =? o4, o1 ↘ o4⟩
  }
  ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 ↘ o3⟩
  List ll = new List();
  ⟨this = o1, x = o2, l = o3, ll = o5 | o1 =? o3, o1 ↘ o3⟩
  ll.n = ll;
  ⟨this = o1, x = o2, l = o3, ll = o5 | o1 =? o3, o1 ↘ o3, o3 n→ o5, ...⟩
  ll.v = x;
}

```

```

public void add(Object x) {
  ⟨this = o1, x = o2 | ε⟩
  List l = this;
  ⟨this = o1, x = o2, l = o1 | ε⟩
  ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 √ o3⟩
  while (l.n != null) {
    ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 √ o3⟩
    l = l.n;
    ⟨this = o1, x = o2, l = o4 | o1 =? o3, o1 √ o3, o3 n→ o4, o1 =? o4, o1 √ o4⟩
    ⟨this = o1, x = o2, l = o4 | o1 =? o4, o1 √ o4⟩
  }
  ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1 √ o3⟩
  List ll = new List();
  ⟨this = o1, x = o2, l = o3, ll = o5 | o1 =? o3, o1 √ o3⟩
  l.n = ll;
  ⟨this = o1, x = o2, l = o3, ll = o5 | o1 =? o3, o1 √ o3, o3 n→ o5, ...⟩
  ll.v = x;
  ⟨this = o1, x = o2, l = o3, ll = o5 | o1 =? o3, o1 √ o3, o3 n→ o5, o5 v→ o2, ...⟩
}

```

```

public void add(Object x) {
  ⟨this = o1, x = o2 | ε⟩
  List l = this;
  ⟨this = o1, x = o2, l = o1 | ε⟩
  ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1  $\swarrow$  o3⟩
  while (l.n != null) {
    ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1  $\swarrow$  o3⟩
    l = l.n;
    ⟨this = o1, x = o2, l = o4 | o1 =? o3, o1  $\swarrow$  o3, o3  $\xrightarrow{n}$  o4, o1 =? o4, o1  $\swarrow$  o4⟩
    ⟨this = o1, x = o2, l = o4 | o1 =? o4, o1  $\swarrow$  o4⟩
  }
  ⟨this = o1, x = o2, l = o3 | o1 =? o3, o1  $\swarrow$  o3⟩
  List ll = new List();
  ⟨this = o1, x = o2, l = o3, ll = o5 | o1 =? o3, o1  $\swarrow$  o3⟩
  l.n = ll;
  ⟨this = o1, x = o2, l = o3, ll = o5 | o1 =? o3, o1  $\swarrow$  o3, o3  $\xrightarrow{n}$  o5, ...⟩
  ll.v = x;
  ⟨this = o1, x = o2, l = o3, ll = o5 | o1 =? o3, o1  $\swarrow$  o3, o3  $\xrightarrow{n}$  o5, o5  $\xrightarrow{v}$  o2, ...⟩
  ⟨... | this =? l, this  $\swarrow$  l, l  $\xrightarrow{n.v}$  x⟩
}

```

```
public void add(Object x) {  
  <... |  $\epsilon$ >  
  List l = this;  
  while (l.n != null) {  
    l = l.n;  
  }  
  List ll = new List();  
  l.n = ll;  
  ll.v = x;  
  <... |  $\text{this} = ? l, \text{this} \setminus l, l \xrightarrow{n.v} x$ >  
}
```

```
public void add(Object x) {  
  <... |  $\epsilon$ >  
  List l = this;  
  while (l.n != null) {  
    l = l.n;  
  }  
  List ll = new List();  
  l.n = ll;  
  ll.v = x;  
  <... |  $\text{this} =^? l, \text{this} \setminus \! \! \! / l, l \xrightarrow{n.v} x$ >  
}
```

- very complex domain – combined may- and must-analysis!

```
public void add(Object x) {  
  <... |  $\varepsilon$ >  
  List l = this;  
  while (l.n != null) {  
    l = l.n;  
  }  
  List ll = new List();  
  l.n = ll;  
  ll.v = x;  
  <... |  $\text{this} =^? l, \text{this} \Downarrow l, l \xrightarrow{n.v} x$ >  
}
```

- very complex domain – combined may- and must-analysis!
- $\Downarrow$  field-insensitive  $\curvearrowright$  post-condition too coarse!

```
public void add(Object x) {  
  <... |  $\epsilon$ >  
  List l = this;  
  while (l.n != null) {  
    l = l.n;  
  }  
  List ll = new List();  
  l.n = ll;  
  ll.v = x;  
  <... |  $\text{this} =^? l, \text{this} \Downarrow l, l \xrightarrow{n.v} x$ >  
}
```

- very complex domain – combined may- and must-analysis!
- $\Downarrow$  field-insensitive  $\curvearrowright$  post-condition too coarse!
- pre-condition?

# A Field-Sensitive May-Analysis

```
public void add(Object x) {  
    List l = this;  
  
    while (l.n != null) {  
        l = l.n;  
    }  
  
    List ll = new List();  
  
    l.n = ll;  
  
    ll.v = x;  
  
}
```



# A Field-Sensitive May-Analysis

```
public void add(Object x) {  
   $\langle \epsilon \rangle$   
  List l = this;  
  
  while (l.n != null) {  
  
    l = l.n;  
  
  }  
  
  List ll = new List();  
  
  l.n = ll;  
  
  ll.v = x;  
  
}
```

# A Field-Sensitive May-Analysis

```
public void add(Object x) {  
   $\langle \epsilon \rangle$   
  List l = this;  
   $\langle \text{this} \xrightarrow{\epsilon} \leftarrow{\epsilon} l \rangle$   
  while (l.n != null) {  
  
    l = l.n;  
  
  }  
  
  List ll = new List();  
  
  l.n = ll;  
  
  ll.v = x;  
  
}
```

# A Field-Sensitive May-Analysis

```
public void add(Object x) {  
   $\langle \epsilon \rangle$   
  List l = this;  
   $\langle \text{this} \xrightarrow{\epsilon} \leftarrow \epsilon l \rangle$   
  while (l.n != null) {  
     $\langle \text{this} \xrightarrow{\epsilon} \leftarrow \epsilon l \rangle$   
    l = l.n;  
  
  }  
  
  List ll = new List();  
  
  l.n = ll;  
  
  ll.v = x;  
  
}
```

# A Field-Sensitive May-Analysis

```
public void add(Object x) {  
   $\langle \epsilon \rangle$   
    List l = this;  
   $\langle \text{this} \xrightarrow{\epsilon} \leftarrow \epsilon l \rangle$   
    while (l.n != null) {  
   $\langle \text{this} \xrightarrow{\epsilon} \leftarrow \epsilon l \rangle$   
      l = l.n;  
   $\langle \text{this} \xrightarrow{n} \leftarrow \epsilon l \rangle$   
    }  
  
    List ll = new List();  
  
    l.n = ll;  
  
    ll.v = x;  
  
}
```

# A Field-Sensitive May-Analysis

```
public void add(Object x) {  
   $\langle \epsilon \rangle$   
  List l = this;  
   $\langle \text{this} \xrightarrow{\epsilon} \leftarrow \langle \epsilon \rangle l \rangle$   $\langle \text{this} \xrightarrow{n?} \leftarrow \langle \epsilon \rangle l \rangle$   
  while (l.n != null) {  
     $\langle \text{this} \xrightarrow{\epsilon} \leftarrow \langle \epsilon \rangle l \rangle$   
    l = l.n;  
     $\langle \text{this} \xrightarrow{n} \leftarrow \langle \epsilon \rangle l \rangle$   
  }  
  
  List ll = new List();  
  
  l.n = ll;  
  
  ll.v = x;  
  
}
```

# A Field-Sensitive May-Analysis

```
public void add(Object x) {  
   $\langle \epsilon \rangle$   
  List l = this;  
   $\langle \text{this} \xrightarrow{\epsilon} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n?} \leftarrow \epsilon l \rangle$   
  while (l.n != null) {  
     $\langle \text{this} \xrightarrow{n?} \leftarrow \epsilon l \rangle$   
    l = l.n;  
     $\langle \text{this} \xrightarrow{n} \leftarrow \epsilon l \rangle$   
  }  
  
  List ll = new List();  
  
  l.n = ll;  
  
  ll.v = x;  
  
}
```

# A Field-Sensitive May-Analysis

```
public void add(Object x) {  
   $\langle \epsilon \rangle$   
  List l = this;  
   $\langle \text{this} \xrightarrow{\epsilon} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n?} \leftarrow \epsilon l \rangle$   
  while (l.n != null) {  
     $\langle \text{this} \xrightarrow{n?} \leftarrow \epsilon l \rangle$   
    l = l.n;  
     $\langle \text{this} \xrightarrow{n.n?} \leftarrow \epsilon l \rangle$   
  }  
  
  List ll = new List();  
  
  l.n = ll;  
  
  ll.v = x;  
  
}
```

# A Field-Sensitive May-Analysis

```
public void add(Object x) {  
   $\langle \epsilon \rangle$   
  List l = this;  
   $\langle \text{this} \xrightarrow{\epsilon} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n?} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n*} \leftarrow \epsilon l \rangle$   
  while (l.n != null) {  
     $\langle \text{this} \xrightarrow{n?} \leftarrow \epsilon l \rangle$   
    l = l.n;  
     $\langle \text{this} \xrightarrow{n.n?} \leftarrow \epsilon l \rangle$   
  }  
  
  List ll = new List();  
  
  l.n = ll;  
  
  ll.v = x;  
  
}
```



# A Field-Sensitive May-Analysis

```
public void add(Object x) {  
   $\langle \epsilon \rangle$   
  List l = this;  
   $\langle \text{this} \xrightarrow{\epsilon} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n?} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n*} \leftarrow \epsilon l \rangle$   
  while (l.n != null) {  
     $\langle \text{this} \xrightarrow{n*} \leftarrow \epsilon l \rangle$   
    l = l.n;  
     $\langle \text{this} \xrightarrow{n.n?} \leftarrow \epsilon l \rangle$   
  }  
  
  List ll = new List();  
  
  l.n = ll;  
  
  ll.v = x;  
  
}
```

# A Field-Sensitive May-Analysis

```
public void add(Object x) {  
   $\langle \epsilon \rangle$   
  List l = this;  
   $\langle \text{this} \xrightarrow{\epsilon} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n?} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n*} \leftarrow \epsilon l \rangle$   
  while (l.n != null) {  
     $\langle \text{this} \xrightarrow{n*} \leftarrow \epsilon l \rangle$   
    l = l.n;  
     $\langle \text{this} \xrightarrow{n.n*} \leftarrow \epsilon l \rangle$   
  }  
  
  List ll = new List();  
  
  l.n = ll;  
  
  ll.v = x;  
  
}
```

# A Field-Sensitive May-Analysis

```
public void add(Object x) {  
   $\langle \epsilon \rangle$   
  List l = this;  
   $\langle \text{this} \xrightarrow{\epsilon} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n?} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l \rangle$   
  while (l.n != null) {  
     $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l \rangle$   
    l = l.n;  
     $\langle \text{this} \xrightarrow{n.n^*} \leftarrow \epsilon l \rangle$   
  }  
   $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l \rangle$   
  List ll = new List();  
  
  l.n = ll;  
  
  ll.v = x;  
  
}
```

# A Field-Sensitive May-Analysis

```
public void add(Object x) {  
   $\langle \epsilon \rangle$   
  List l = this;  
   $\langle \text{this} \xrightarrow{\epsilon} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n?} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n*} \leftarrow \epsilon l \rangle$   
  while (l.n != null) {  
     $\langle \text{this} \xrightarrow{n*} \leftarrow \epsilon l \rangle$   
    l = l.n;  
     $\langle \text{this} \xrightarrow{n.n*} \leftarrow \epsilon l \rangle$   
  }  
   $\langle \text{this} \xrightarrow{n*} \leftarrow \epsilon l \rangle$   
  List ll = new List();  
   $\langle \text{this} \xrightarrow{n*} \leftarrow \epsilon l \rangle$   
  l.n = ll;  
  
  ll.v = x;  
  
}
```

# A Field-Sensitive May-Analysis

```
public void add(Object x) {  
   $\langle \epsilon \rangle$   
  List l = this;  
   $\langle \text{this} \xrightarrow{\epsilon} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n?} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l \rangle$   
  while (l.n != null) {  
     $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l \rangle$   
    l = l.n;  
     $\langle \text{this} \xrightarrow{n.n^*} \leftarrow \epsilon l \rangle$   
  }  
   $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l \rangle$   
  List ll = new List();  
   $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l \rangle$   
  l.n = ll;  
   $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l, l \xrightarrow{n} \leftarrow \epsilon ll \rangle$   
  ll.v = x;  
}
```

# A Field-Sensitive May-Analysis

```
public void add(Object x) {  
   $\langle \epsilon \rangle$   
  List l = this;  
   $\langle \text{this} \xrightarrow{\epsilon} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n?} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l \rangle$   
  while (l.n != null) {  
     $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l \rangle$   
    l = l.n;  
     $\langle \text{this} \xrightarrow{n.n^*} \leftarrow \epsilon l \rangle$   
  }  
   $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l \rangle$   
  List ll = new List();  
   $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l \rangle$   
  l.n = ll;  
   $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l, l \xrightarrow{n} \leftarrow \epsilon ll \rangle$   
  ll.v = x;  
   $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l, l \xrightarrow{n} \leftarrow \epsilon ll, ll \xrightarrow{v} \leftarrow \epsilon x \rangle$   
}
```

# A Field-Sensitive May-Analysis

```
public void add(Object x) {  
   $\langle \epsilon \rangle$   
  List l = this;  
   $\langle \text{this} \xrightarrow{\epsilon} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n?} \leftarrow \epsilon l \rangle$   $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l \rangle$   
  while (l.n != null) {  
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   $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l, l \xrightarrow{n} \leftarrow \epsilon ll \rangle$   
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   $\langle \text{this} \xrightarrow{n^*} \leftarrow \epsilon l, l \xrightarrow{n} \leftarrow \epsilon ll, ll \xrightarrow{v} \leftarrow \epsilon x \rangle$   
   $\langle \text{this} \xrightarrow{n^+.v} \leftarrow \epsilon x \rangle$   
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- very complex domain – combined may- and must-analysis!
- $\nabla$  field-insensitive  $\curvearrowright$  post-condition too coarse!
- pre-condition?



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  - if  $\langle \epsilon \rangle$  holds before add,  $\langle \text{this} \xrightarrow{n^+.v} \leftarrow^{\epsilon} x \rangle$  holds after add
  - all side-effects of add are captured by  $\langle \text{this} \xrightarrow{n^+.v} \leftarrow^{\epsilon} x \rangle$

- implementation with support for...

# Current State of Development

- implementation with support for...
  - non-recursive Java programs

- implementation with support for...
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  - acyclic data-structures



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# Current State of Development

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  - ~~loop unrolling~~
- TODO
  - cyclic data-structures



- novel heap shape analysis

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## Demo!

## Analyzed property: Are all data-structures trees/DAGs?

300 examples from TPDB

283×same result

8×HashMap

tree vs. DAG	1
tree vs. arbitrary	1
DAG vs. arbitrary	6
arbitrary vs. tree	1

Table: ~2 weeks ago

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DAG vs. arbitrary	4
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Table: yesterday

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  - 1 year of development

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## Note

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  - 9 years of development
  - highly optimized for TPDB
  - infers “binary” results
- new approach
  - 1 year of development
  - unoptimized
  - infers detailed results